Coverless Text Information Hiding Method Using the Frequent Words Hash

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Abstract

The attackers may discover the existence of the secret information or even get it by analyzing the cover's statistical characteristics, changes of which often occur due to the embedding. In this paper, a novel coverless text information hiding method was proposed. By using the words rank map and the frequent words hash, normal texts containing the secret information could be retrieved from the text database, and will be sent to the receiver without any modification. Because the embedding is not needed, the proposed method could be able to escape from almost all state-of-the-art steganalysis methods.

Keywords: Big Data; Coverless Information Hiding; Frequent Words Hash; Rank Map; Steganography

1 Introduction

Steganography, also known as information hiding, is a secure communication method that conveys secret messages in the form of plaintexts so that the appearances of the secret messages will not draw eavesdroppers' attention while they are being transmitted through an open channel [17]. It can be used for intellectual property protection and secret communications [9]. For example, reference [20, 25] introduced two methods of detecting illegal copies of copyrighted images. For information hiding, there are many kinds of covers, such as texts [15], images [7, 18], videos [16], etc. [1, 12, 21].

Compared to the image or other covers, text information hiding is the most difficult kind of steganography due to the lack of redundancy. Because the text is frequently used in people's daily lives, however, text information hiding has attracted many researchers' interest, and has many results [23]. Classified by the covers, text steganography could be put into three types: text format-

based [3, 10, 11], generating-based and embedding-based natural language information hiding. For text formatbased information hiding, the embedded information will no longer exist if the document is generated without format after extracting the text content. Generating-based natural language information hiding methods can fool the computer statistical analysis, but is relatively easy to be identified by people [2]. Embedding-based natural language information hiding methods have more robust and better concealment than text format-based information hiding, but the hiding algorithm is difficult to implement, and there are some deviations and distortions in the statistic and linguistics because of the limitation of the natural language processing [13].

Once the information hiding algorithm is public, the steganalysis methods will be appearing. So attackers will know the existence of the secret information by analyzing changes of statistical characteristics of the covers caused by the embedded information. Be there an algorithm with which the secret information could be hidden without any modification of the covers. Coverless information hiding [24, 6], firstly proposed by Xingming Sun et al., is the best answer to the above question. Reference [24] presented a coverless image steganography framework, and Reference [6] proposed a coverless text information hiding method. These two methods can directly retrieve the stego-image (stego-text) without any modification of the covers. Recently, coverless information hiding, which requires no modification on covers and could resist various steganalysis technologies, draws more and more attention from researchers [5, 22].

In this paper, a novel coverless text information hiding method is proposed. Firstly, a text database is constructed by collecting a large number of texts from the Internet. Then the word rank maps of the words will be calculated by statistically analyzing the text big data, meanwhile, the frequent words distance of every text is the words' occurrence in a text database (or in a text), the calculated. When a certain information will be transmitted, a normal text containing the secret information is retrieved from the text database by using the frequent words distance and the word rank maps, and sent to the receiver without any modification.

$\mathbf{2}$ **Coverless Information Hiding**

Coverless information hiding is a new challenging research field. In fact, "coverless" is not to say that there is no carrier, but compared with the conventional information hiding, coverless information hiding requires no other carries [6]. The idea of coverless information hiding is often used in our daily life, and the acrostic poem is a classic example. An acrostic poem is shown in Figure 1 form which we can learn that the secret information is "TREE". Coverless information hiding is essentially the disclosure of secret information in the text. Its distinctive characteristic is "no embedding", that is, a carrier cannot embed secret information by modifying it [6].



Figure 1: An acrostic poem

3 The Proposed Method

3.1Preparation of the Text Database

We construct a natural text database by fetching the news from the normal news web sites. For each word of the vocabulary, we calculate the frequency of its occurrence, and then rank the words with the descending way (Most frequent word has rank 1, next frequent word has rank 2 ...). Figure 2 shows the ranking result of words in a text database. In order to make good use of the information of

Word Rank Map of a text database (or a text) is defined as:

$$RM = \{(w_i, f_i) | i = 1, 2, \cdots, U\}$$
(1)

where U is the number of unique words in a text database (or in a text), i is the rank of a word w_i , and f_i is the frequency of w_i . Figure 2 shows the word rank map of a text database. For the example in Figure 2, we can obtain

$$RM = \{(the, 124020), (and, 55654), \\ (of, 54550), (to, 52331), \cdots \}$$
(2)

Obviously, the top frequent words are: the, and, of, to, in, a, on, for, and etc.

total	number of	words:	1967696	
total	number of	unique	words: 5763	8
rank	word	1	frequency:	
1	the		124020 :	
5	and		55654	
â -	of		54550	
Ā	to		\$2331	
÷ .	in		46181	
6		:	41390	
2	on.		19241	
	for	:	17883	
ä	that		17338	
ĩo	said		17230	
11	i.e.		16107	
12	china		13051	
12	with	:	13542	
14	a.c.	:	11784	
12	hu		10478	
16	by	:	10336	
19	all.		10230	
16			9900 1	
18	wi i i		0100	
12	ne		9199	
20	Trom		9079	
21	nas		8960	
22	was		0001	
23	5		7703	
62	De		7389 :	
22	nave	-	7251 :	
20	are		7145 :	
27	an .		7011 :	
28	chinese	-	6992 :	
29	115	:	6119 :	
30	015		5587 :	

Figure 2: Part of a text database word rank map

For each text in a text database, we can obtain its word rank map defined as Equation (1). Figure 3 shows the rank map of a text named as "2.1 million Audi cars affected by emissions cheating scandal.txt", in which there are 185 words, 113 unique words. From the rank map, we learn that the top frequent words are: the, in, software, emission, cars, etc.

For the top frequent words in the text database, we can calculate their occurrences in a text in the same collection. So, the Frequent Words Hash Function is defined as:

$$H_k(t) = \{h_1 h_2 h_3 \cdots h_k\}$$
(3)

where k is the number of the top frequent words chose form the vocabulary of a text database, t is a text in the text database, and h_i is defined by:

$$h_i = \begin{cases} 1 & \text{the i-th frequent word} \\ & \text{appears in text t} \\ 0 & \text{the i-th frequent word does not} \\ & \text{appear in text t} \end{cases}$$
(4)

total	number of words number of uniqu	: 185 e words: 113
rank	word :	frequency:
1	the :	19 :
2	in :	6 :
3	software :	4 :
4	emission :	4 :
s	cars :	4 :
6	car i	4 :
7	of :	4 :
8	on :	4 :
9	a :	4 :
10	audi :	3 :
11	to :	3 1
12	epa :	3 :
13	cheating :	3 :
14	million :	3 :
15	by :	3 :
16	said :	2 :
17	1 :	2 :
18	total :	2 :
19	that :	2 :
20	were :	2 :
21	agency :	2 1
22	protection:	2 :
23	emissions :	2 :
24	environmental	: 2:
25	scandal :	2 :
26	vehicles :	2 :
27	5 :	2 :
28	u :	2 :
29	volkswagen:	2 :
30	COBDADY 1	2 .

Figure 3: Part of a text's word rank map

For the top 30 frequent words shown in Figure 2, we can calculate the hash value of a text named as "2.1 million Audi cars affected by emissions cheating scandal.txt". The hash value is:

$$H_k(t) = \{111111101110111101001010010010\}$$
(5)

So, we map a text into a 30 bits string. Figure 4 shows the hash values of some texts in a text database. In order to measure the occurrence of the frequent words in a text t, we define the Frequent Words Distance of a text as:

$$DFW_k(t) = HD(H_k(t), (b_1, b_2, \cdots, b_k))$$

$$b_i = 0, \quad i = 1, 2, \cdots, k.$$
(6)

where k is the number of the top frequent words, t is a text in the text database, and HD is Hamming Distance calculating operation.

111111	111	0.1.1	1.1.1.	1001	0100	010110
1111111	111	101	001	0001	0001	100000
1111111	111	001	110	1111	0101	001001
1101111	100	001	010	0100	0001	001001
1111111	111	101	011	1011	0101	001001
1111111	101	101	111	1011	1111	101011
1111111	111	101	111	1111	1101	101010
1111111	111	001	111	0101	1100	111000
1111111	111	101	111	1111	1111	111010
1111111	111	101	111	1101	1101	110010
1111111	111	101	110	1 1 0 1	1101	110010
1111111	111	100	011	0111	1101	000000
1011111	011	000	011	1100	1001	001010
1111111	111	101	111	1111	1101	110011
1111111	1 1 1	101	111	1110	1101	011010
1111111	101	101	010	1100	0000	101010
1111111	1 1 1	011	110	1001	1000	111110
1111111	110	101	011	1111	0100	101001
1111111	1 1 1	101	011	1100	1101	111010
1111111	100	001	110	0001	1100	011010
1111111	1 1 1	101	111	1111	1101	111011
1111111	111	111	110	0101	0101	010110
1111111	1 1 1	100	111	1101	0011	111010
1111111	100	001	011	1000	0110	100001
1111111	111	101	111	1011	1111	100001
1111111	110	101	111	1111	1111	110001
1111111	000	101	011	0101	1001	010001
1111111	111	101	111	1111	1111	111011

Figure 4: The hash values of some texts in a text database

By statistical analyzing the occurrence of each word in the text database, we can calculate the word rank map words in W as the right part of Figure 7. Suppose the

of each word. For each word w_i appearing in the text database, its word rank map is defined as

$$RMW_i = \{(rw_{ij}, fw_{ij}, wt_{ij}) | i = 1, 2, \cdots, U; \\ j = 1, 2, \cdots, N\}$$
(7)

where rw_{ij} is the rank of the word w_i in a text wt_{ij} according to its occurrence, fw_{ij} is the frequency of w_i 's occurrence in wt_{ij} , and N is the number of the texts in which w_i appears. Figure 5 shows the word rank map of "from", whose rank is 20 in the word rank map as shown in the Figure 2.

010-	11 Off	
ank-	-Freque	encySource Text
6	3	'Containing China' a Japanese strateg
2	4	'Cyber Monday' sales set to hit recor
0	3	'Disturbance' only harms China, Japan Fu Yin
2	2	'Hobbit' vanished far earlier than researcher:
2	1	'In' leads by 7 percentage points ahead of Br
	7	'Maternity tourism' in US sees chang
3	5	'Mini Messi', 5, aims to follow in star's foo
5	4	'No more survivors' in collapsed overpas
1	2	'Plane wreckage' found in Thailand fuels talk
2	3	'Slave' women held for 30 yrs rescue
1	2	'Symbol of rebirth' opens in NY
51	1	1 killed, 51 injured in Mexico candy factory
0	3	1 killed, dozens wounded at UN base in S. Sud.
9	2	1 survived, 53 bodies recovered in Algeria pla
.6	3	10 Afghan militants killed in army operatio
3	4	10 facts about HIVAIDS on World AIDS Da
2	2	10 killed in flash flood in eastern Indi
33	1	10 US Navy sailors detained by Iran repor
13	1	100 killed in tribal clashes in Sudan's Darfu
9	2	11 trapped miners rescued from S. African min
0	2	12 China-made helicopters delivered to Cambod
4	2	12 IS militants killed in air strike in weste
0	2	125 killed in month-long battles of Irag's Anl
8	4	126 hostages rescued, 4 attackers killed as Br
	4	126 rescued, 4 attackers killed as operation i
7	1	12th Elephant Festival held in Nepa

Figure 5: Part of the rank map of "from"

3.2**Information Hiding**

The information hiding process is shown in Figure 6. Detail procedures are introduced as follows. Suppose the constructed text database is T, and the communication key is k. We can calculate the word rank map of T by using Equation (1), and get the vocabulary of T, and let it be $W = \{w_i | i = 1, 2, \dots, U\}$ where i is the rank of w_i , and U is the number of unique words in T. For each text t_i in T, we can calculate its word rank map by using Equation (1), and let it be RM_{t_i} .



Figure 6: The process of information hiding

Because the key is k, we arrange the top frequent $k \times k$

hidden message is $M = m_1, m_2, \cdots, m_n$ where m_i is a word, and n is the number of words in the hidden message. For each word m_i in M, it is chosen from the top frequent $k \times k$ words in W. Obviously, the selection range of m_i depends on k. Therefore, both sides of communication **Finally**, as described above, we can search a stego-text can choose more k so that m_i has more options.

w_1	Ħ	W ₁	w ₂	W3		w_k
<i>w</i> ₂	∉	$= w_{k+1}$	W_{k+2}	w_{k+3}		w_{2k}
:	Ť	_ :	:	1	:	:
W_{k-1}	1∉	$W_{(k-2)\times k+1}$	$W_{(k-2)\times k+2}$	$W_{(k-2)\times k+3}$		$W_{(k-1)\times k}$
Wk	⇐	$W_{(k-1)\times k+1}$	$W_{(k-1)\times k+2}$	$W_{(k-1)\times k+3}$		$W_{k \times k}$

Figure 7: The words conversion table

Words Conversion 3.2.1

In order to enhance the security of the secret message, we convert each word in M into one of the top frequent k words in W before the information hiding. The conversion rule is shown in Figure 7. For each word in $\{w_1, w_2, \cdots, w_k\}$, it will be converted into w_1 . For each word in $\{w_{(k+1)}, w_{(k+2)}, \cdots, w_k\}$, it will be converted into w_2 . And so on.

For each word m_i in M, we can get its rank by using the word rank map of T, and let it be R_{m_i} . Then, it is located in the $((R_{m_i} - 1)/k + 1)$ row, $((R_{m_i} - 1)\% k + 1)$ column in the word conversion table shown in Figure 7, where "%" is a remainder operation. Therefore it will be converted into $m'_i = w_{((R_{m_i}-1)/k+1)}$. In this way, we can convert the secret information $M = m_1, m_2, \cdots, m_n$ into $M' = m'_1, m'_2, \cdots, m'_n$, and M' is a subset of the k top frequent words in W.

3.2.2Searching the Stego-text

For each word $m'_i = w_{((R_{m_i}-1)/k+1)}$ in M', the stego-text is get as follows:

- **Firstly**, for each text t in the text database, we calculate the hash value of $H_k(t)$, defined as Equation (3). where k is the communication key. Then, we can get the frequent words distance of t by using Equation (6), ant let it be $DFW_k(t)$.
- Secondly, because m'_i is in the top k frequent words in the word rank map of T, we can get the rank map of m'_i by using Equation (7), and let it be $RMW_{m'_i}$.
- **Thirdly,** by using the word rank map $RMW_{m'_i}$, we retrieve all texts containing m'_i to search a text t in which the rank of m'_i is equal to $DFW_k(t)$, and the frequency of m'_i 's occurrence is equal to the secret message $M = m_1, m_2, \cdots, m_n$.

 $((R_{m_i}-1)\% k+1)$. There may be some texts satisfying this condition, then, we can select a text from those texts as the stego-text for m'_i .

set for each m'_i in M'. These stego-texts is a normal text set that contains the converted secret message, and they can be sent to the receiver without any modifying.

Information Extraction 3.3

The process of extraction is shown in Figure 8. Suppose the stego-text is S, so S is a set of normal texts. The number of texts in S is the number of words in secret message M. Let k be the communication key. Because the text database T is open for all users, receiver can calculate the word rank map of T, and get the top kfrequent words in W by using the communication key k. Certainly, receiver can get the same word conversion table shown in Figure 7. For each stego-text t in S, the details of information extraction will be introduced as follows.



Figure 8: Secret message extracting process

3.3.1Get the Candidate Word

Receiver can calculate the word rank map of t by using Equation (1), and the frequent words distance of t by using Equation (6) and let it be $DFW_k(t)$. By retrieving the word rank map of t, receiver can get the candidate word whose rank is equal to $DFW_k(t)$ in text t. Obviously, the candidate word is $m'_i = w_{((R_m, -1)/k+1)}$.

Get the Secret Message 3.3.2

By using the word rank map of t, receiver can find the word frequency of m'_i in t, and let it be Fm'_i . Obviously, Fm'_i is equal to $((R_{m_i}-1)\% k+1)$. Receiver can find the secret message m_i that is located in the " $w_{((R_{m_i}-1)/k+1)}$ " row, the $((R_{m_i}-1)\% k+1)$ column in the word conversion table shown in Figure 7.

So, receiver can get every word m_i in M. and then get

4 Discussion

4.1 An Example

In order to clearly explain the above coverless text information hiding process, we illustrate it by a simple example. We have constructed a text database which can be expanded constantly, and it is open for all users. Suppose the communication key is 30 and the secret information M is "mutual visit". It is worth mentioning that, however, M is a subset of the top 900 frequent words in the text database. Both sides of communication may choose larger k, so that words of M have more options. The operating procedure of information hiding is introduced as follows:

Firstly, sender computes the rank map of the text database and ones of each text in it. Because the communication key is 30, sender can obtain the top 30 frequent words set $W_{top30} = \{w_i | i = 1, 2, \dots, 30\}$, and the W_{top30} is: $\{the, and, of, to, in, a, on, for, that, said, is, China, with, as, by, at, it, will, he, form, has, was, s, be, have, are, an, Chinese, its, his <math>\}$.

So, sender can get the word conversion table shown in Figure 9. By retrieving the word rank map of text database, sender finds the rank of "mutual" is 605, and the rank of "visit" is 183. According to the word conversion table shown in Figure 9, therefore, "mutual" is located in 21st row, 5th column, and "visit" is located in 7th row, 3rd column. Hence, "mutual" will be converted into "has", and "visit" will be converted into "on".

the 🖨	$- w_1$	w_2	w_3		w_{30}
and 🖛	w ₃₁	w ₃₂	w ₃₃		W ₆₀
: +	- :	:	:	:	:
its 🖛	- w ₈₄₁	w_{842}	w_{843}		w ₈₇₀
his 🖨	■ W ₈₇₁	W872	w ₈₇₃		W900

Figure 9: The word conversion table when key is 30

Secondly, sender calculates the word rank map of "has" and ones of "on". By retrieving the two word rank maps, sender can find a text named "25reuters-golfryder-usa-north.txt", and let it be t_1 , and a text named "155_Chinese_loggers'_release_not_victory_of _diplomacy.txt", and let it be t_2 . Their word rank maps are shown in Figure 10 and Figure 11. From Figure 10, we learn that the rank of "has" is 24 and its frequency is 5 in text t_1 whose frequent words distance is 24, and the frequent words are $\{the, a, and, of, in, on, at, to, with, has, was, is, that, as, his, for, be, an, it, by, will, he, s, have \}$ in it. From Figure 11, we learn that the rank of "on" is 25 and its frequency is 3 in text t_2 whose frequent words distance is 25, and the frequent words are {the, and, of, to, in, a, on, for, that, is, China, with, as, by, it, will, he, from, has, be, have, are, an, Chinese, its} in it.

2 Denotery golf cycler was north lot i Notepad Ris Edit Fannat View Help					
rank	:word	:	frequency:		
12	the a	1	26 :		
3	and he's		16		
67	of appearanc	es:	13 :		
8 9	in previous		11 1		
11	at		9		
13 14 15	with cup ryder		8 7 7		
16	2014 really		7		
18 19 20	great		6 :		
21	played		6		
24	has team		5		
26 27	was 2010	1	5 :		

Figure 10: The word rank map of a stego-text

Finally, sender sends the two texts t_1 , t_2 as the stegotexts to the receiver.

213,Ohneie	higger: printer, solution, sta	fightern any	old - Notepad
File Edit For	mat Van Help		
rank	:word	:	frequency:
1	the	:	26 :
2	of	:	18 :
3	chinese	:	15 :
4	myanmar	:	12 :
5	in	:	11 :
6	should	:	10 :
7	they	:	10 :
8	and	:	9 :
9	to	:	8 :
10	a		8 :
11	not	:	7 :
12	were	:	7 :
13	be		7 :
14	are		6 :
15	for		6 :
16	155		<u>Š</u> :
17	as	-	5 :
18	these		5
19	have	-	Ś :
20	released		4
21	by		4
22	īŧ	-	4
23	sentenced		4
24	we		4
25	on	-	3
26	loggers		3
27	is	-	3
28	has		3 :
29	that		
30	illegal	:	3 :

Figure 11: The word rank map of a stego-text

Because the text database is open to all users, receiver can calculate its word rank map, the top 30 frequent words and the word conversion table shown in Figure 9 by using the communication key k = 30. Then, he (or she) calculates the frequent words distance of t_1 and ones of t_2 , and finds that they are 24 and 25. So, receiver retrieves the word rank maps of t_1 and t_2 , and gets the candidate words "has" and "on" whose ranks are 24 and 25. From the word rank map of the stego-texts, he (or she) also learns that the candidate words' frequency are 5 and 3. Finally, receiver gets the secret message "mutual" located in "has" row, 5th column in the word conversion table shown in Figure 9, and "visit" located in "on" row, 3rd column. Hence, receiver successfully extract the secret message "mutual visit" from the stego-texts.

4.2 Security Analysis

Steganalysis is usually performed through the use of irrelevance between the embedded information and the carriers. Attackers often make steganalysis by analyzing the difference of their statistical distributions [19]. In our proposed hiding method, however, the carriers are normal pure text and the secret information is not been embedded in the carriers. The carriers can be sent to receiver without any modification. So the information hiding does not change the probability distribution of the carriers. According to the definition of the security of an information hiding system in [4], the proposed information hiding method is theoretically safe. At the same time, the proposed approach is also followed the Kerckhoffs Principle [8] in cryptography, and detail of information hiding is open. If he does not know the communication key, the attacker cannot gain any information about the hidden information [14]. Therefore, the proposed method could resist almost all kinds of current steganalysis method.

4.3 The Importance of Big Data

However, it is worth mentioning that, in order to enhance security, there are two works must be done: one is to change periodically the communication key to ensure that the secret message may be converted into different subsets of the top frequent words in the text database. The second is to establish a large text database (text big data) to increase the probability of the frequent words distance is equal to the rank of a word in a text, and so there are more choices of the stego-texts [23]. For example, in the chose text database, for the word "China", there are 3258 texts in which it appears, and their frequent words distance are shown in Figure 12. From the Figure 12, we learn that these values are not evenly distributed. There are 37 texts whose frequent words distance is 12, and there is only one text in which the rank of "China" is 12. Therefore, the text big data is necessary to ensure the smoothly implement of the proposed method.

Because the text big data is an important guarantee of the smooth implementation of the proposed method, some files should reside in the memory buffer when the big data is handling. We firstly calculate the word rank map of each text in the text big data, then the word rank map of the text database, and finally the ones of each word of the vocabulary, so the computing cost is expensive especially when computing the word rank map of the text database. In order to reduce the complexity, we will use the "inverted index" for storage optimization.



Figure 12: The distribution of frequent words distances of texts containing "China"

Because the location lab is simply designed, the capacity of the proposed method is one word per text. In order to increase the capacity of information hiding, we will design better lab location methods in the future.

5 Conclusion

This paper presented a coverless text information hiding method based on the frequent words hash. By using the words rank map and the frequent words hash, normal texts containing the secret information could be retrieved from the text database, and will be sent to the receiver without any modification. Because there is no embedding, the information hiding does not change the probability distribution of the covers. Therefore, the proposed method is theoretically safe, and could be able to escape from almost all state-of-the-art steganalysis methods.

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